

```
#include <stdlib.h>
#include <string.h>
#include <ctype.h>

#define MAXPAROLA 30
#define MAXRIGA 80

int main(int argc, char *argv[])
{
    int freq[MAXPAROLA]; /* vettore di contatori
delle frequenze delle lunghezze delle parole */
    char riga[MAXRIGA];
    int i, inizio, lunghezza;
    FILE *f;

    for(i=0; i<MAXPAROLA; i++)
        freq[i]=0;

    if(argc != 2)
    {
        fprintf(stderr, "ERRORE: serve un parametro con il nome del file\n");
        exit(1);
    }
    f = fopen(argv[1], "r");
    if(f==NULL)
    {
        fprintf(stderr, "ERRORE: impossibile aprire il file %s\n", argv[1]);
        exit(1);
    }

    while( fgets( riga, MAXRIGA, f ) != NULL )
```



Graphs

Single Source Shortest Paths for DAGs

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Shortest paths on weighted DAGs

- ❖ For a DAG the SSSPs problem can be solved with a simplified algorithm
- ❖ Shortest paths are always well defined even if there are negative-weight edges
 - This is because, obviously, negative-weight cycles cannot exist in a DAG

Shortest paths on weighted DAGs

❖ As there are no cycles it is enough to

➤ Topologically sort the DAG

- Impose a linear order on the vertices

Perform a DFS computing
end-processing times
Order vertices using the end-
processing times

➤ Relax all vertices following the sorted order given
by the topological sort

- In other words, it suffices to make just one pass
over the vertices in the topological sorted order
- As we process a vertex, we relax each edge that
leaves the vertex

SSSP for DAGs

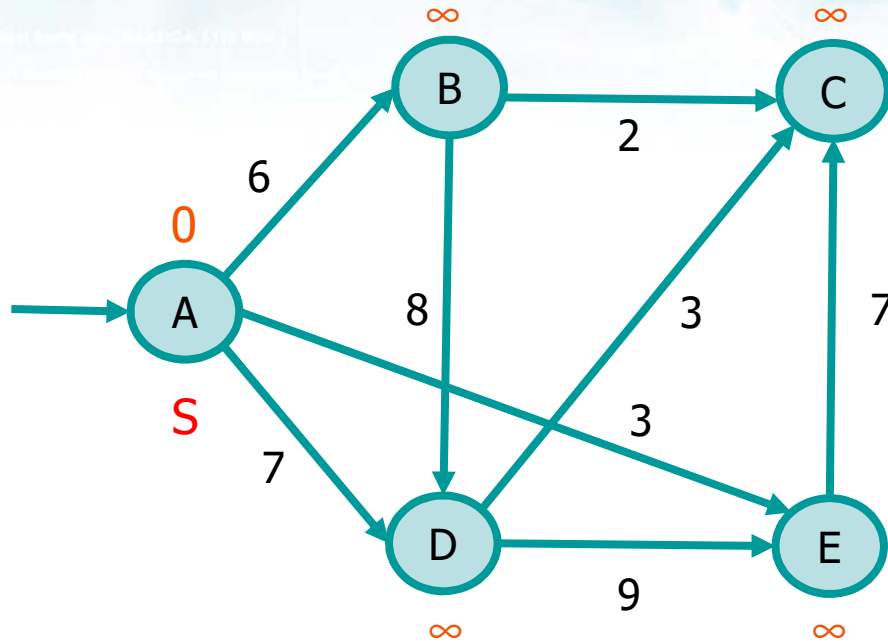
Pseudo-code

```
sssp_for_DAGs (G, w, s)
  topological sort the vertices of G
  initialize_single_source (G, s)

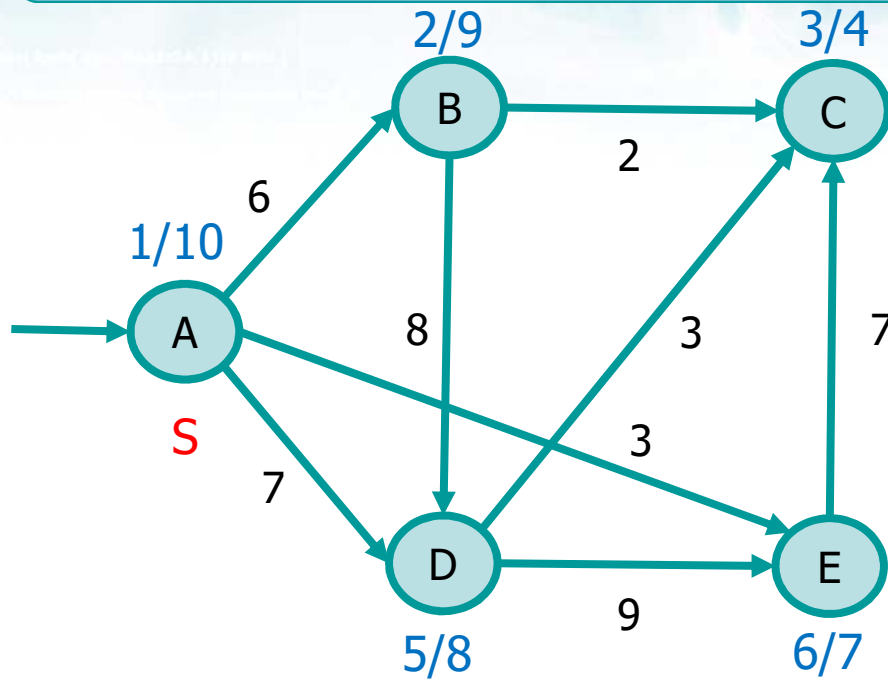
  for each vertex u ∈ V
    for each vertex v ∈ adjacency list of u
      relax (u, v, w)
```

Taken in topologically sorted order

Example

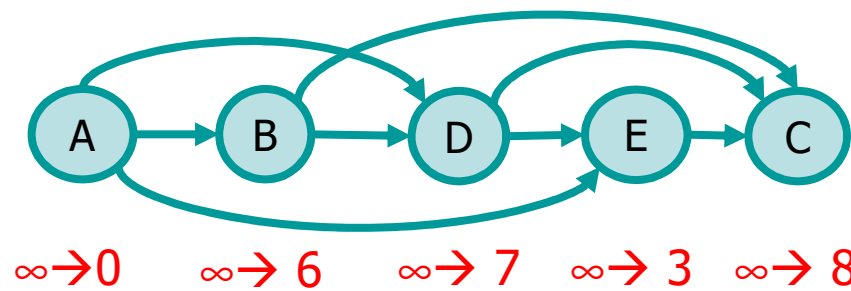


Solution



Relaxation order

(A, B)
(A, D)
(A, E)
(B, D)
(B, C)
(D, E)
(D, C)
(E, C)



Complexity

Pseudo-code

```

sssp_for_DAGs (G, w, s)
  topological sort the vertices of G
  initialize_single_source (G, s)

  for each vertex u ∈ V
    for each vertex v ∈ adjacency list of u
      relax (u, v, w)
    
```

$\Theta(|V| + |E|)$

$\Theta(|V|)$

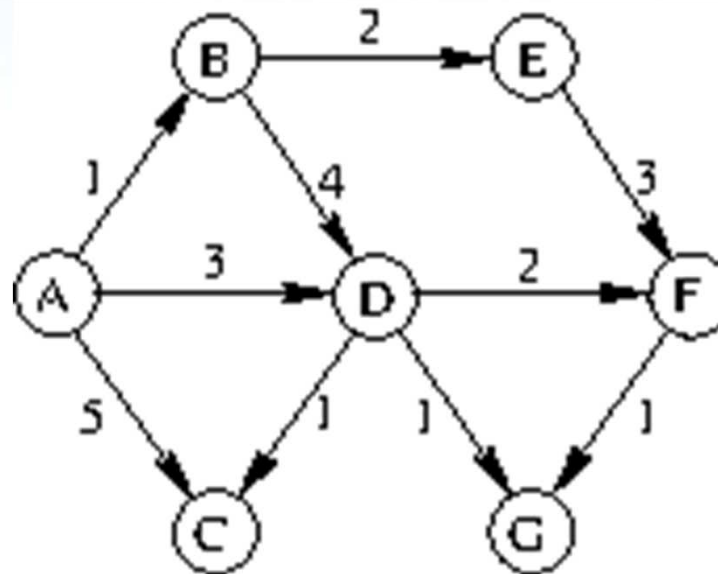
Executed E times
alltogether

$\Theta(1) \rightarrow \Theta(|E|)$

Taken in topological
sorted order

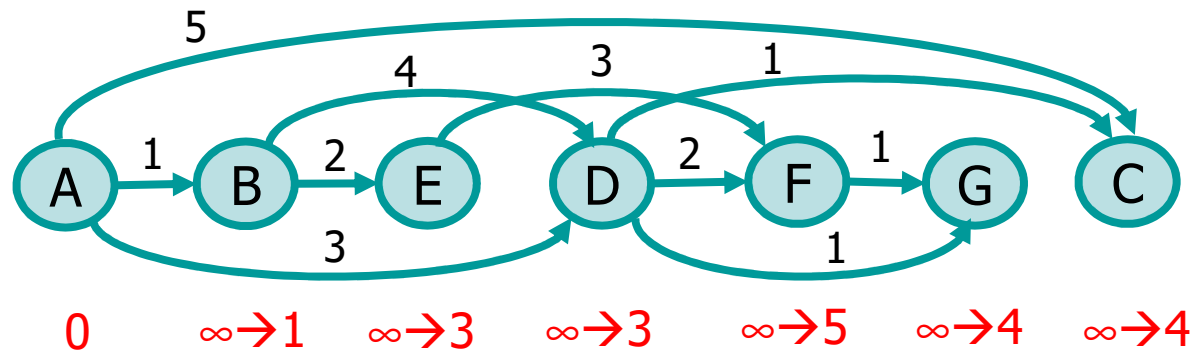
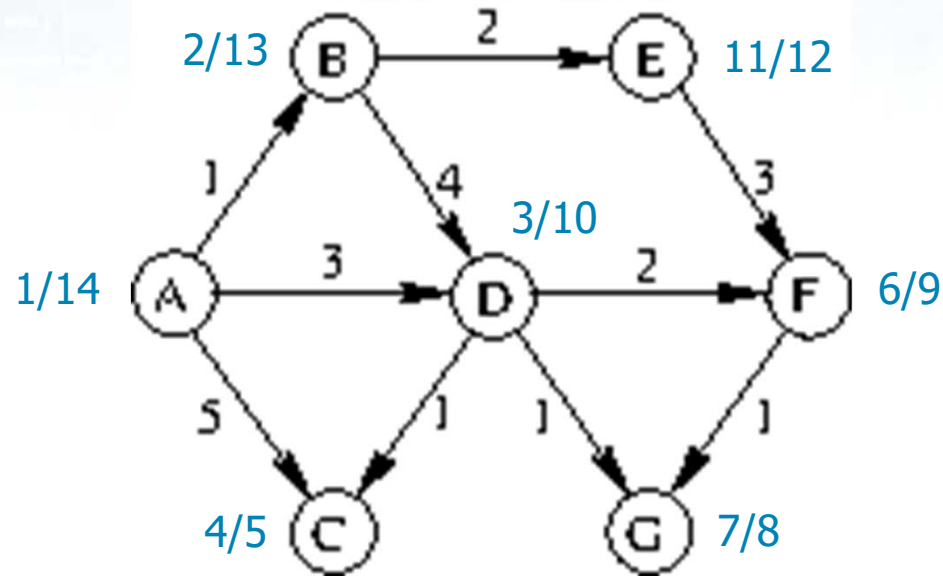
Overall running time complexity
 $T(n) = \Theta(|V| + |E|)$

Exercise

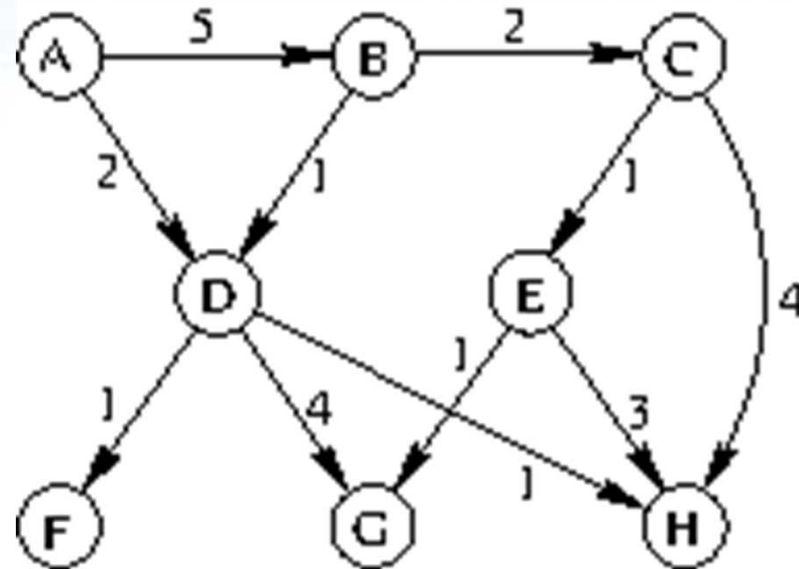


- ❖ Given the following graph find all shortest paths starting from vertex A

Solution

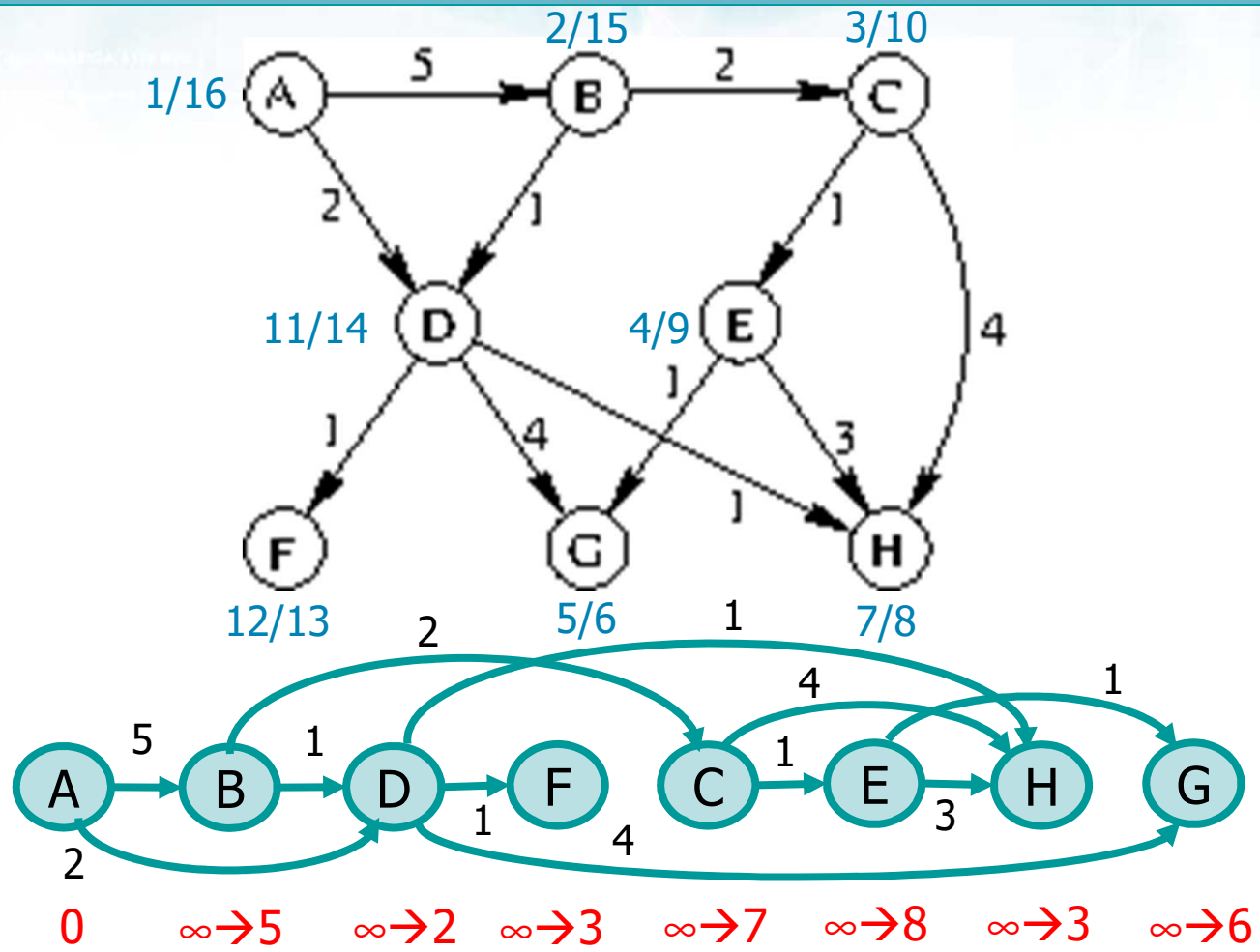


Exercise



- ❖ Given the following graph find all shortest paths starting from vertex A

Solution

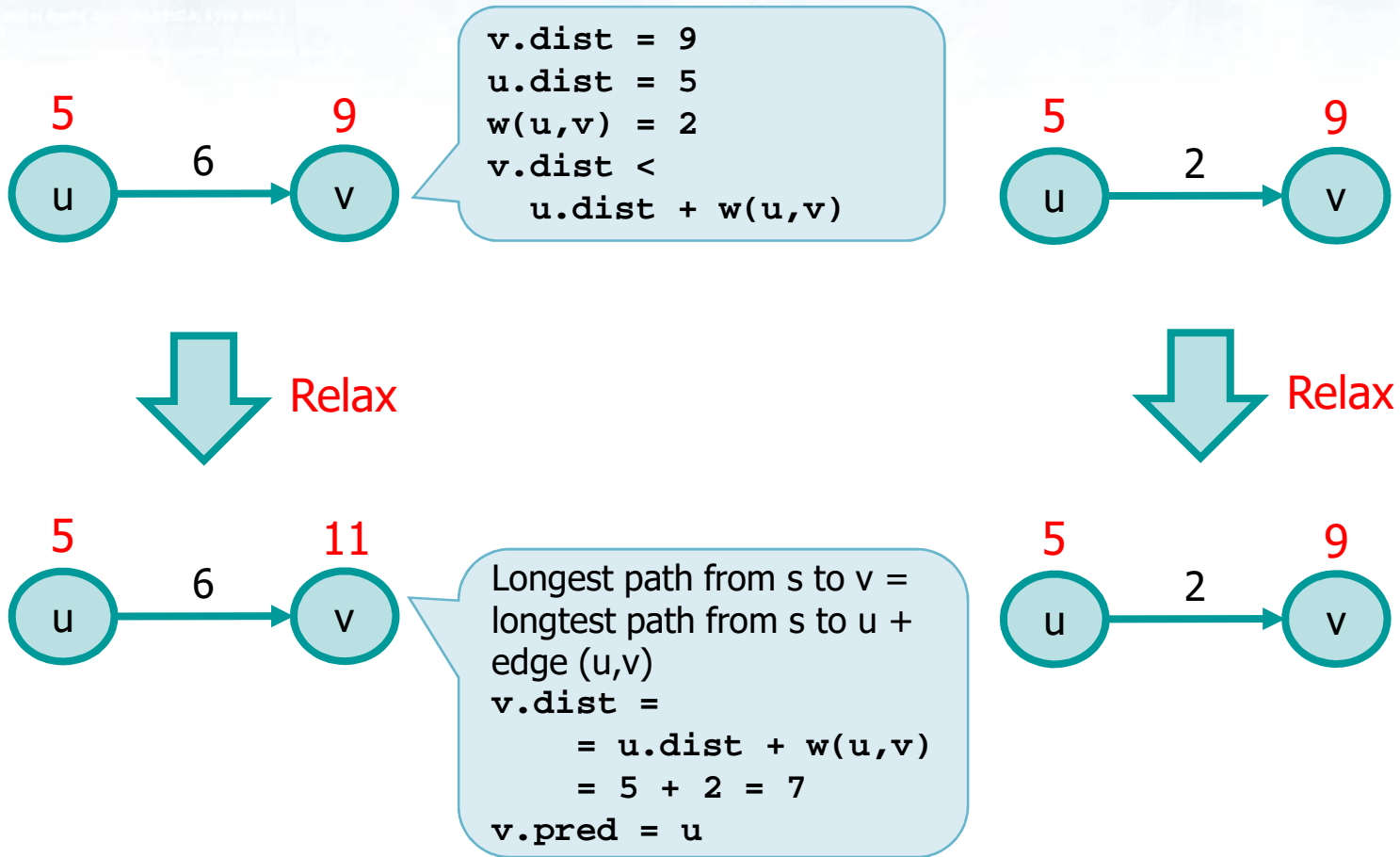


Longest path on weighted DAG

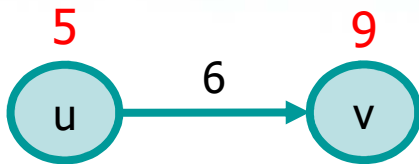
- ❖ Problem intractable on generic weighted graph
- ❖ As on a DAG there are no cycles, the problem become computationally feasible
 - Topologically sort the DAG
 - For all ordered vertices
 - Apply the "inverse" relaxation rule starting from that vertex

```
inverse_relax (u, v, w) {  
  if (v.dist < u.dist + w(v,u)) {  
    v.dist = u.dist + w(v,u)  
    v.pred = u  
  }  
}
```

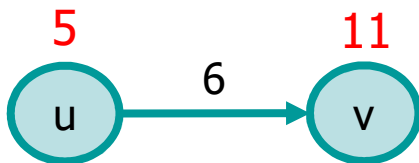
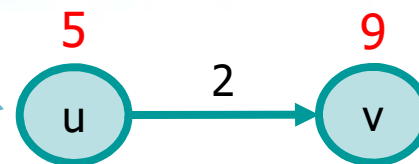
Example



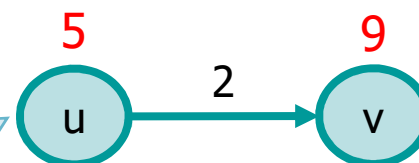
Example



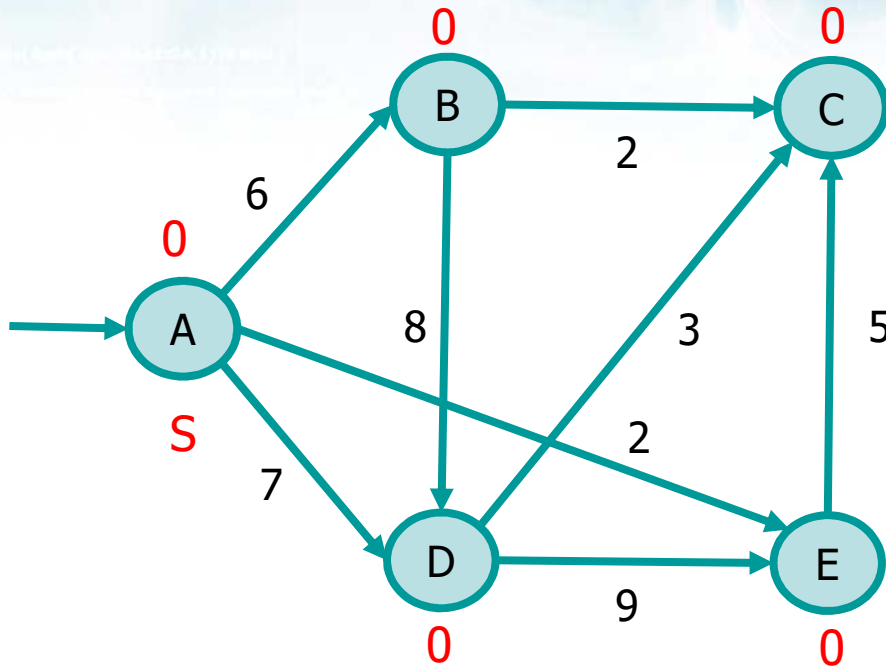
`v.dist = 9`
`u.dist = 5`
`w(u,v) = 2`
`v.dist >`
`u.dist + w(u,v)`



Relaxation has no effect
`v.dist = unchanged`
`= 6`
`v.pred = unchanged`

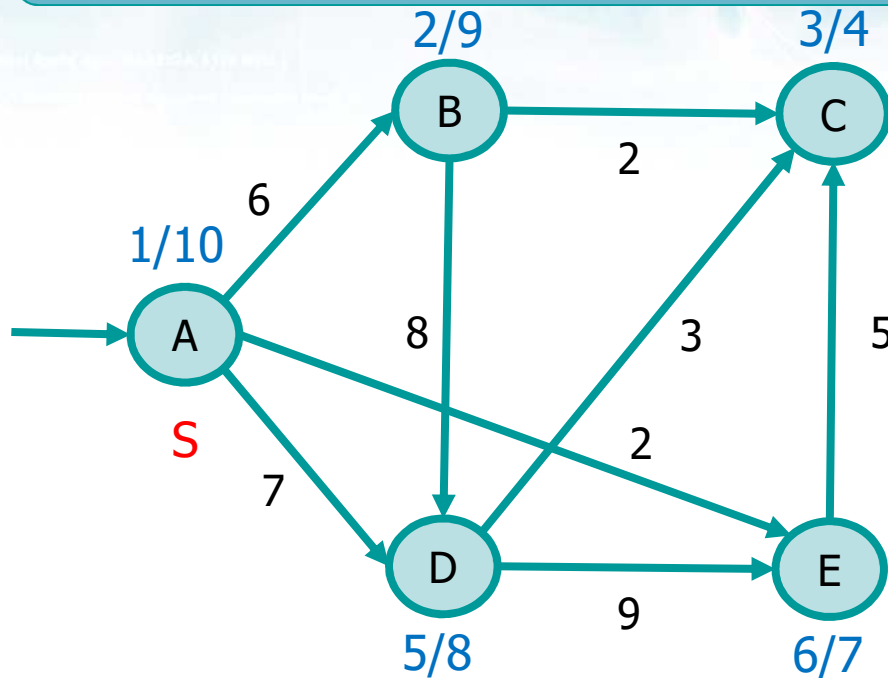


Example



The initial estimate is equal to zero for all vertices

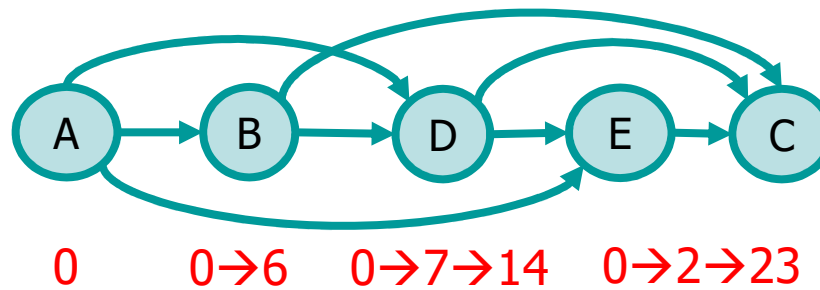
Solution



Relaxation order

- (A, B)
- (A, D)
- (A, E)
- (B, D)
- (B, C)
- (D, E)
- (D, C)
- (E, C)

Relaxation order

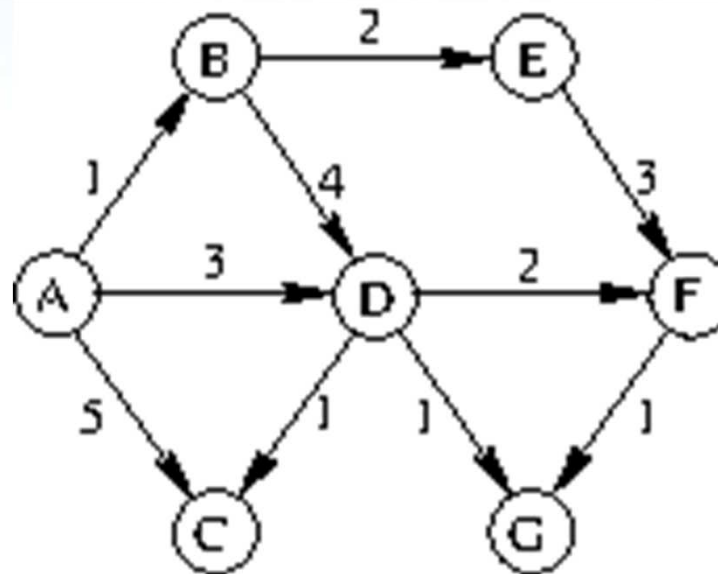


0 → 8 → 17 → 28

Complexity

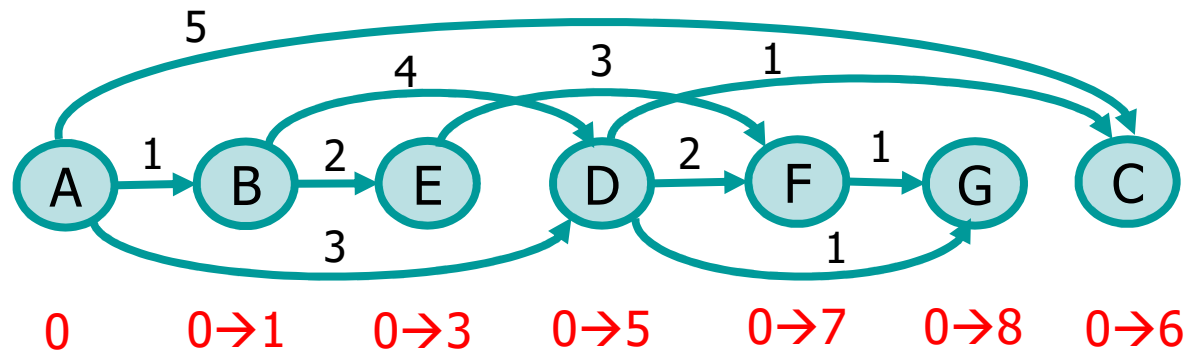
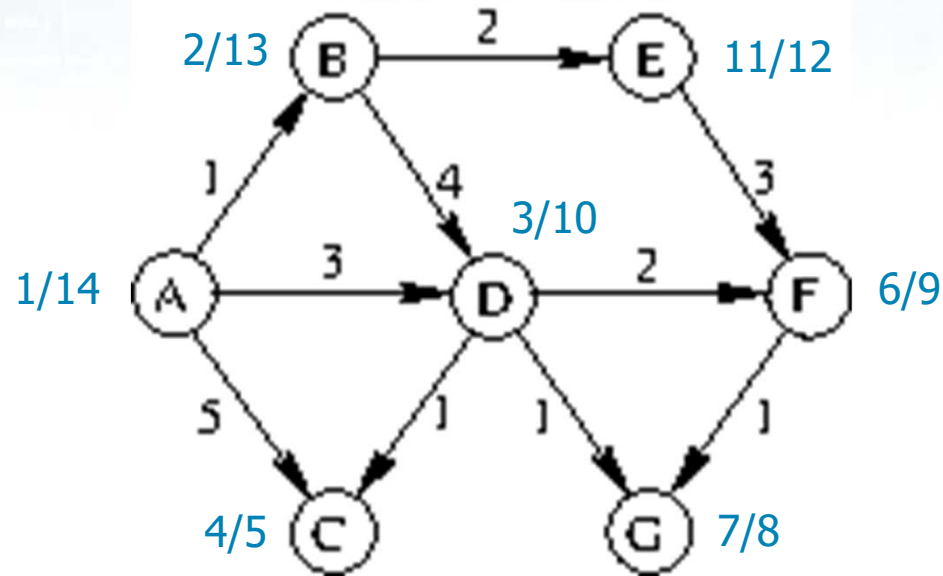
- ❖ As the algorithm for the shortest paths for DAGs

Exercise

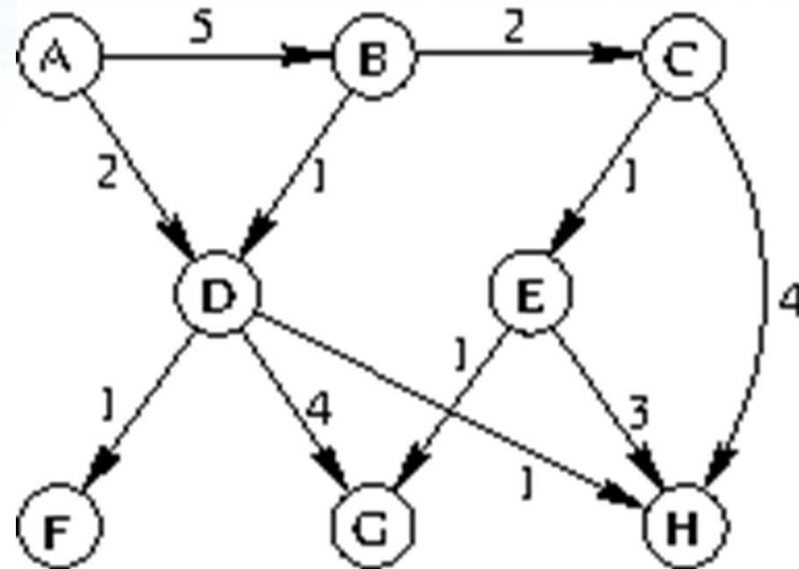


- ❖ Given the following graph find all longest paths starting from vertex A

Solution



Exercise



- ❖ Given the following graph find all longest paths starting from vertex A

Solution

