

```
#include <stdlib.h>
#include <string.h>
#include <ctype.h>

#define MAXPAROLA 30
#define MAXRIGA 80

int main(int argc, char *argv[])
{
    int freq[MAXPAROLA]; /* vettore di contatori
delle frequenze delle lunghezze delle parole */
    char riga[MAXRIGA];
    int i, inizio, lunghezza;
    FILE *f;

    for(i=0; i<MAXPAROLA; i++)
        freq[i]=0;

    if(argc != 2)
    {
        fprintf(stderr, "ERRORE: serve un parametro con il nome del file\n");
        exit(1);
    }
    f = fopen(argv[1], "r");
    if(f==NULL)
    {
        fprintf(stderr, "ERRORE: impossibile aprire il file %s\n", argv[1]);
        exit(1);
    }

    while( fgets( riga, MAXRIGA, f ) != NULL )
```

Trees and BSTs

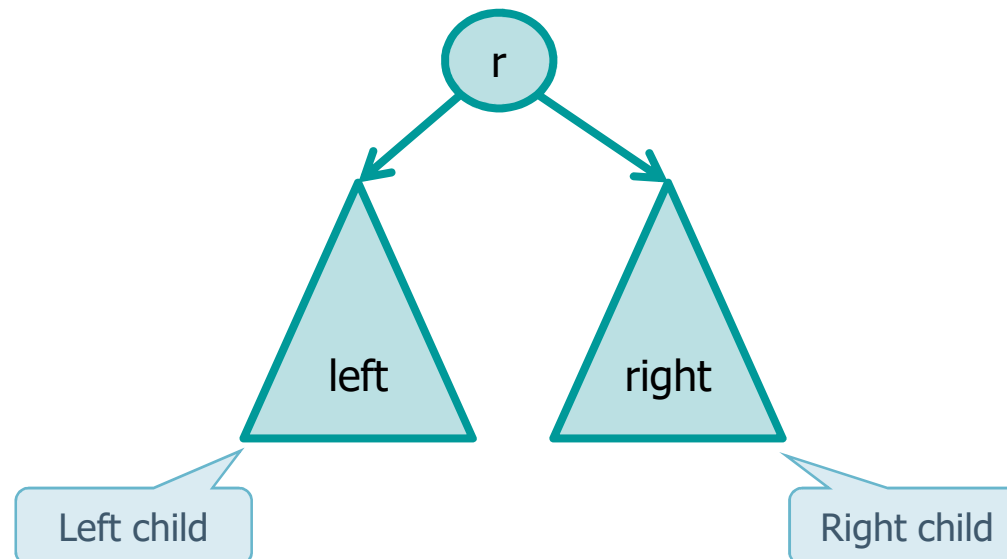
Trees

Paolo Camurati and Stefano Quer
Dipartimento di Automatica e Informatica
Politecnico di Torino

Binary Trees

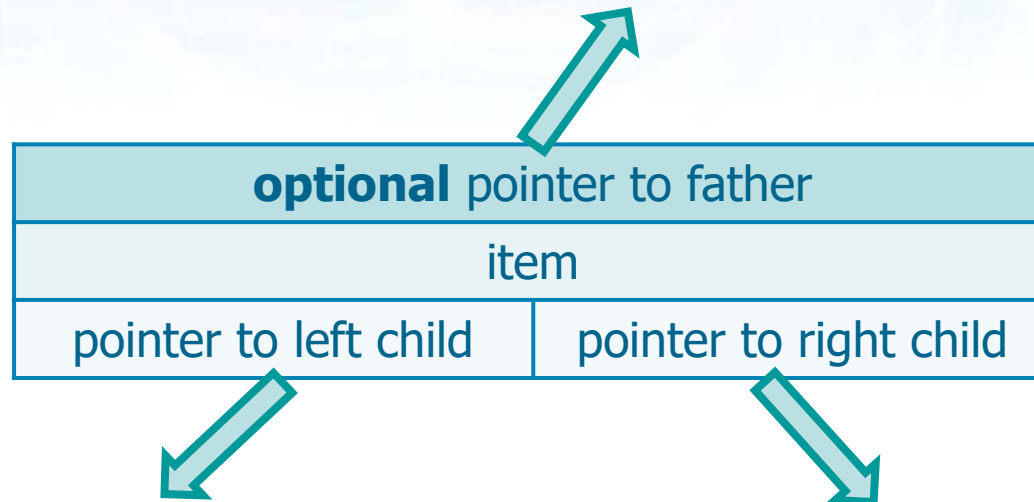
❖ Recursive definition

- Empty set of nodes
- Root, left subtree, right subtree



Binary Trees

item → key
is a string
(in this section)



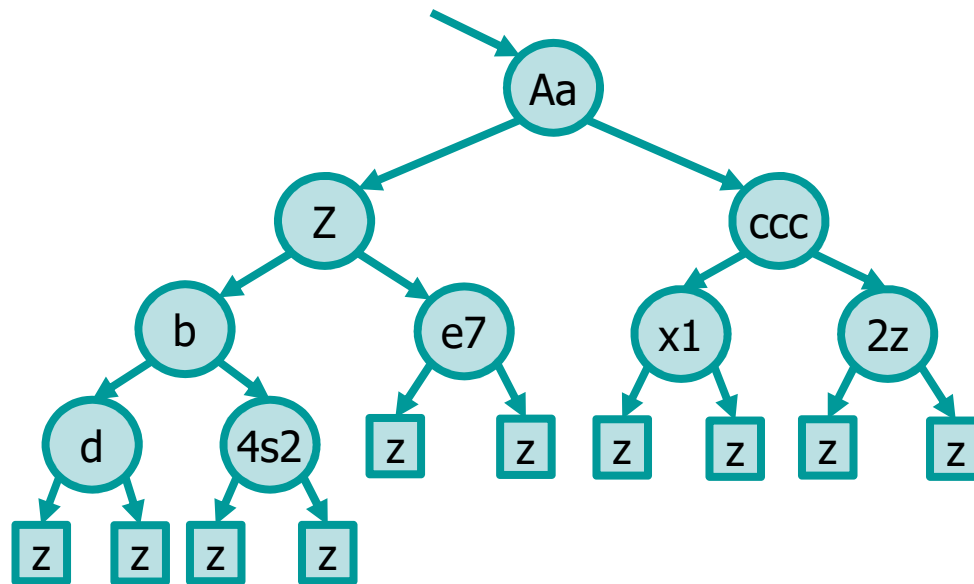
```
typedef struct node *link;  
struct node {  
    Item item;  
    link l;  
    link r;  
};
```

ADT: We use functions
to compare keys, etc.

Binary Trees

❖ Tree

- Access through pointer to root

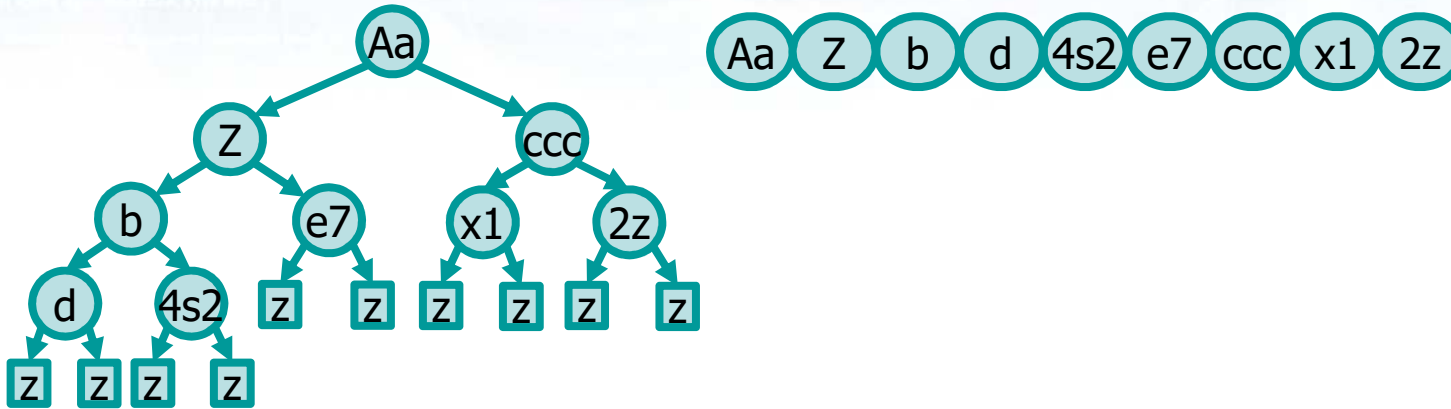


- Dummy sentinel node z or NULL pointer

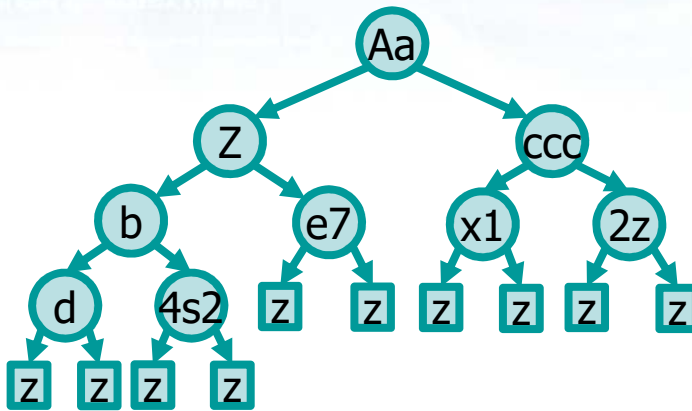
Visits

- ❖ A tree traversal or a tree visit lists the nodes according to a strategy
- ❖ Three strategies are generally used
 - Pre-order
 - **Root**, Left child (l), Right child (r)
 - In-order
 - Left child (l), **Root**, Right child (r)
 - Post-order
 - Left child (l), Right child (r), **Root**

Pre-order



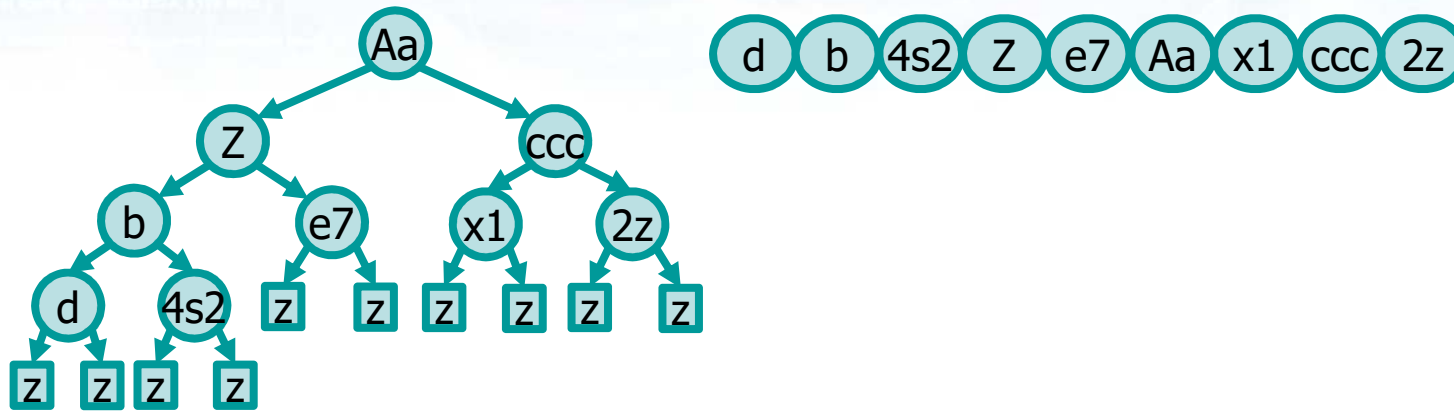
Pre-order



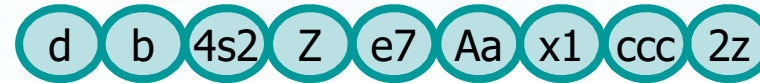
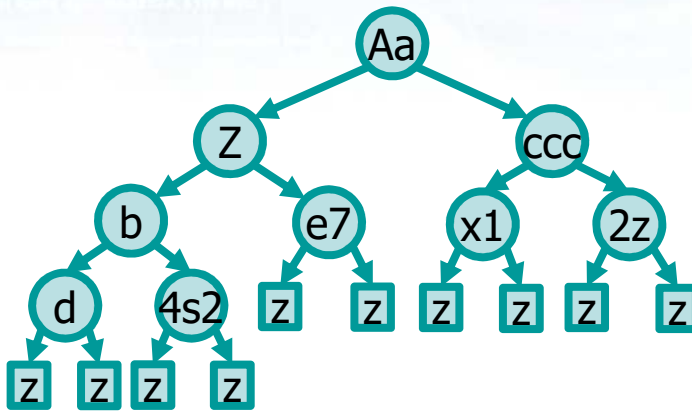
Function **item_print** writes keys

```
void preorder_r (
    link root,
    link z
) {
    if (root == z)
        return;
    item_print (root->item);
    preorder_r (root->l, z);
    preorder_r (root->r, z);
    return;
}
```

In-order

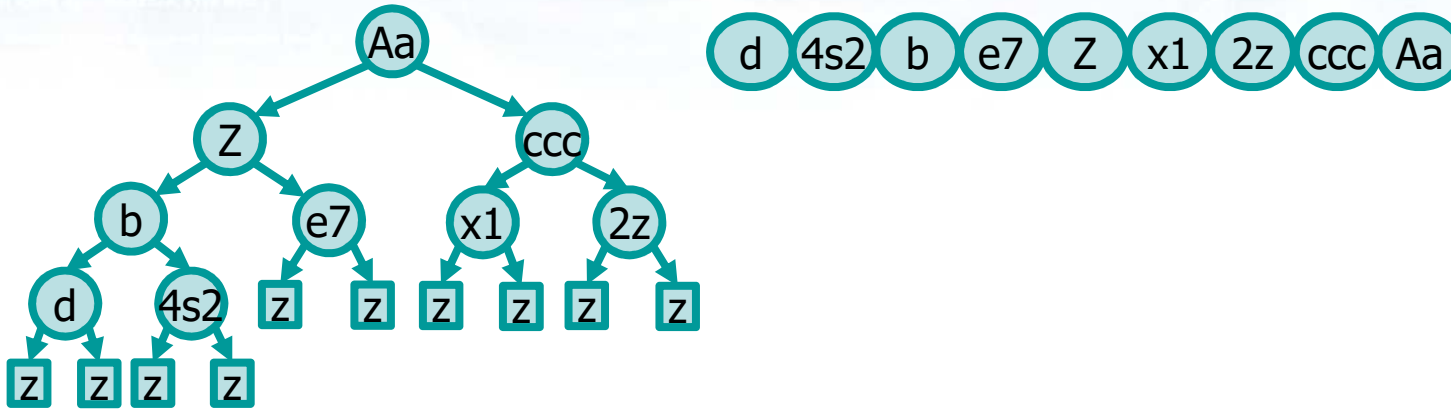


In-order

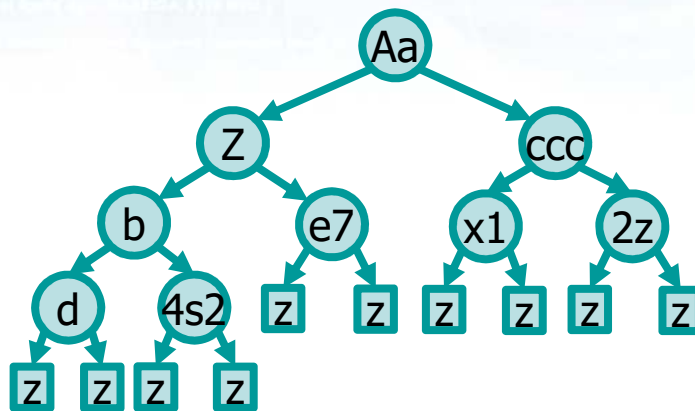


```
void inorder_r (
    link root,
    link z
){
    if (root == z)
        return;
    inorder_r (root->l, z);
    item_print (root->item);
    inorder_r (root->r, z);
    return;
}
```

Post-order



Post-order



```
void postorder_r(
    link root,
    link z
){
    if (root == z)
        return;
    postorder_r (root->l, z);
    postorder_r (root->r, z);
    item_print (root->item);
    return;
}
```

Comparison

```
void preorder_r (  
    link root,  
    link z  
) {  
    if (root == z)  
        return;  
    item_print (root->item);  
    preorder_r (root->l, z);  
    preorder_r (root->r, z);  
    return;  
}
```

```
void inorder_r (  
    link root,  
    link z  
) {  
    if (root == z)  
        return;  
    inorder_r (root->l, z);  
    item_print (root->item);  
    inorder_r (root->r, z);  
    return;  
}
```

```
void postorder_r(  
    link root,  
    link z  
) {  
    if (root == z)  
        return;  
    postorder_r (root->l, z);  
    postorder_r (root->r, z);  
    item_print (root->item);  
    return;  
}
```

Complexity Analysis

❖ Case 1

➤ Complete tree

- $D(n) = \Theta(1)$
- $C(n) = \Theta(1)$
- $a = 2$
 - Two sub-problems
- $b = 2$
 - Overall size (of the two partitions) $n-1$, conservatively approximated to n , i.e., $n/2$ and $n/2$

Divide
and
conquer

➤ Recurrence equation

- $T(n) = 1 + 2 \cdot T(n/2) \quad n > 1$
- $T(1) = 1 \quad n = 1$

```
void inorder_r (...) {  
    if (root == z)  
        return;  
    inorder_r (root->l, z);  
    item_print (root->item);  
    inorder_r (root->r, z);  
    return;  
}
```

Complexity Analysis

➤ With unfolding

- $T(n) = 1 + 2 \cdot T(n/2)$
- $T(n/2) = 1 + 2 \cdot T(n/4)$
- $T(n/4) = 1 + 2 \cdot T(n/8)$
- ...
- $T(1) = 1$

```
void inorder_r (...) {
    if (root == z)
        return;
    inorder_r (root->l, z);
    item_print (root->item);
    inorder_r (root->r, z);
    return;
}
```

Termination condition

$$\frac{n}{2^i} = 1$$

$$i = \log_2 n$$

➤ That is

- $T(n) = 1 + 2 \cdot (1 + 2 \cdot T(n/4))$
- $= 1 + 2 \cdot (1 + 2 \cdot (1 + 2 \cdot T(n/8))) = 1 + 2 + 4 \dots$
- $= \dots$
- $= \sum_{i=0}^{\log n} 2^i = \frac{(2^{\log n + 1} - 1)}{2 - 1} = 2 \cdot 2^{\log n} - 1 = 2n - 1$
- $= \mathbf{O(n)}$

Complexity Analysis

❖ Case 2

➤ **Totally unbalanced tree**

➤ The tree degenerates into a list

- $D(n) = \Theta(1)$
- $C(n) = \Theta(1)$
- $a = 1$
- $k_i = 1$

➤ Recurrence equation

- $T(n) = 1 + T(n-1) \quad n > 1$
- $T(1) = 1 \quad n = 1$

```
void inorder_r (...) {  
    if (root == z)  
        return;  
    inorder_r (root->l, z);  
    item_print (root->item);  
    inorder_r (root->r, z);  
    return;  
}
```

Complexity Analysis

➤ With unfolding

- $T(n) = 1 + T(n-1)$
- $T(n-1) = 1 + T(n-2)$
- $T(n-2) = 1 + T(n-3)$
- ...
- $T(1) = 1$

➤ That is

- $T(n) = 1 + 1 + T(n-2)$
 $= 1 + 1 + 1 + T(n-3)$
 $= \dots$
 $= \sum_{i=1}^n 1 = n = \mathbf{O(n)}$

```
void inorder_r (...) {  
    if (root == z)  
        return;  
    inorder_r (root->l, z);  
    item_print (root->item);  
    inorder_r (root->r, z);  
    return;  
}
```


Parameter Computation

- ❖ Compute the number of nodes of a binary tree

Number of nodes

Root and Sentinel (or nothing
for a termination condition
checking on NULL)

```
int count (link root, link z) {  
    int l, r;  
  
    if (root == z)  
        return 0;  
  
    l = count (root->l, z);  
    r = count (root->r, z);  
    return (l+r+1);  
}
```

Parameter Computation

- ❖ Compute the height of a binary tree

```
int height (link root, link z) {
    int u, v;

    if (root == z)
        return -1;

    u = height (root->l, z);
    v = height (root->r, z);

    if (u>v)
        return (u+1);
    else
        return (v+1);
}
```

Exercise

- ❖ Given an n-ary tree compute its
 - Number of nodes
 - Height

Node definition

```
typedef struct node *link;
struct node {
    Item item;
    int degree;
    link *child;
};
```

Array of size degree
Eventually
link child[DEGREE];

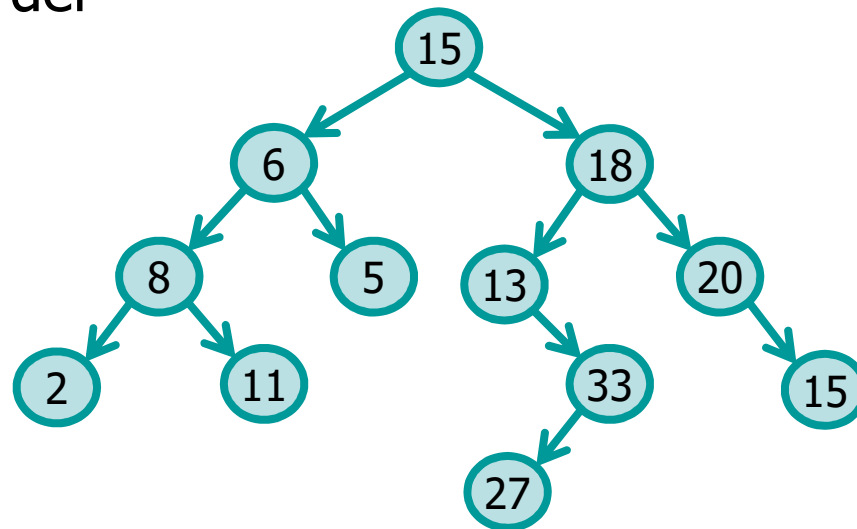
Solution

```
int count (link root, link z) {
    int i, c;
    if (root == z)
        return 0;
    for (c=0, i=0; i<root->degree; i++) {
        c = c + count (root->child[i], z);
    }
    return (c+1);
}

int height (link root, link z) {
    int i, tmp, max=-1;
    if (root == z)
        return -1;
    for (i=0; i<root->degree; i++) {
        tmp = height (root->child[i], z);
        if (tmp > max)
            max = tmp;
    }
    return (max+1);
}
```

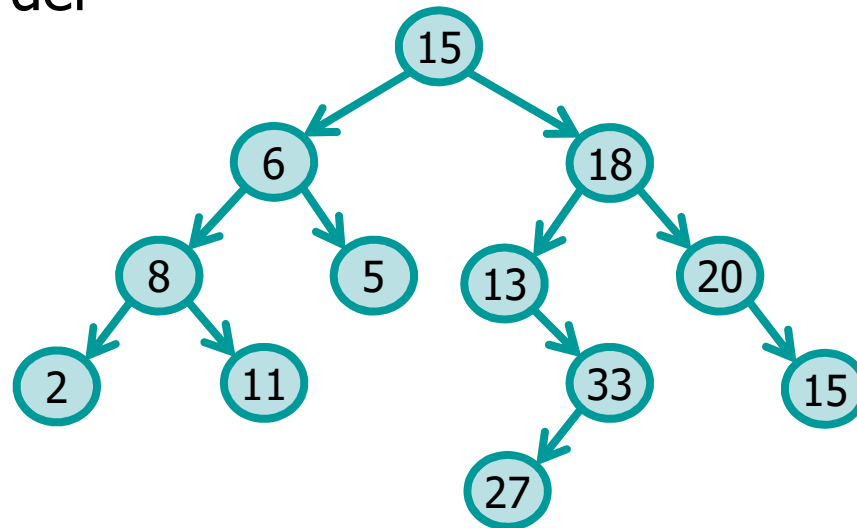
Exercise

- ❖ Given the following tree visit it in pre, in, and post-order



Solution

- ❖ Given the following tree visit it in pre, in, and post-order



- Pre-order : 15 6 8 2 11 5 18 13 33 27 20 15
- In-order : 2 8 11 6 5 15 13 27 33 18 20 15
- Post-order: 2 11 8 5 6 27 33 13 15 20 18 15

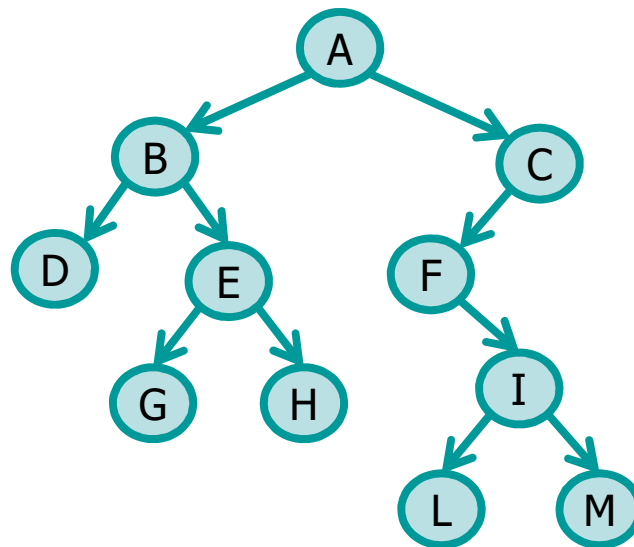
Exam: 29 January 2018

Exercise

- ❖ Consider a binary tree with 11 nodes
- ❖ Draw it considering that its pre, in and post-order visits return the following sequences
 - Pre-order: A B D E G H C F I L M
 - In-order: D B G E H A F L I M C
 - Post-order: D G H E B L M I F C A

Solution

- Pre-order: A B D E G H C F I L M
- In-order: D B G E H A F L I M C
- Post-order: D G H E B L M I F C A



Application: Expressions

- ❖ Given an algebraic expression (brackets to change operator priority), it is possible to build the corresponding tree according to the simplified grammar

```
<exp> = <operand> | <exp> <op> <exp>  
<operand> = A .. Z  
<op> = + | * | - | /
```

Termination
Condition

Recursion

Example

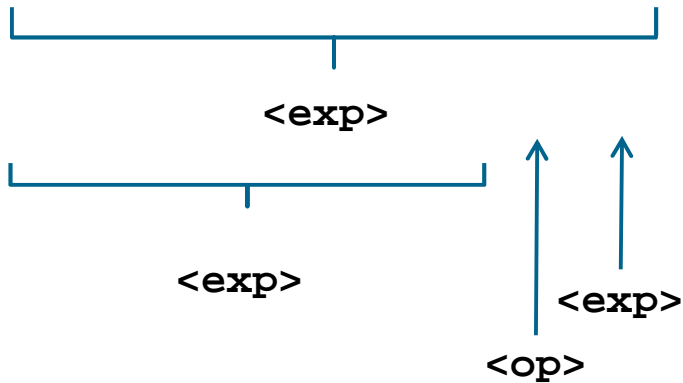
- ❖ Using the following grammar (left-hand side) parse the following equation (right-hand side)

```
<exp> = <operand> | <exp> <op> <exp>  
<operand> = A .. Z  
<op> = + | * | - | /
```

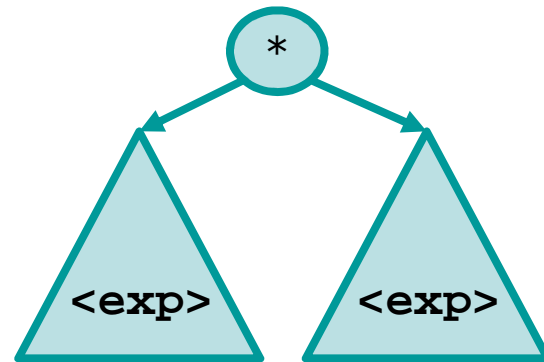
$[(A + B) * (C - D)] * E$

Solution: Step 1

$[(A + B) * (C - D)] * E$



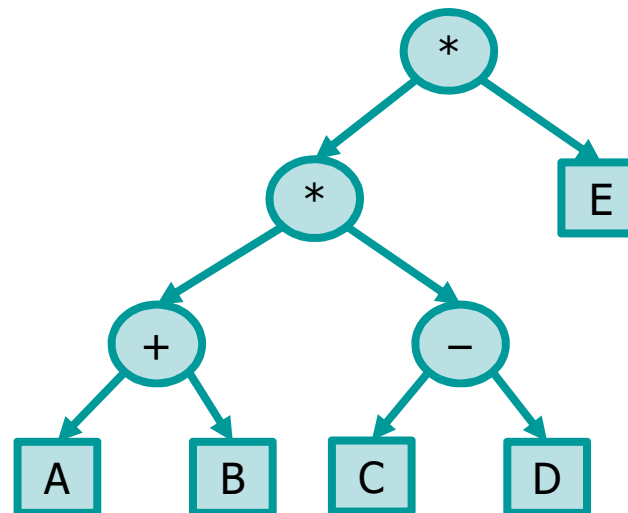
$\langle \text{exp} \rangle = \langle \text{operand} \rangle \mid \langle \text{exp} \rangle \langle \text{op} \rangle \langle \text{exp} \rangle$
 $\langle \text{operand} \rangle = A \dots Z$
 $\langle \text{op} \rangle = + \mid * \mid - \mid /$



Solution

$[(A + B) * (C - D)] * E$

$\langle \text{exp} \rangle = \langle \text{operand} \rangle \mid \langle \text{exp} \rangle \langle \text{op} \rangle \langle \text{exp} \rangle$
 $\langle \text{operand} \rangle = A \dots Z$
 $\langle \text{op} \rangle = + \mid * \mid - \mid /$



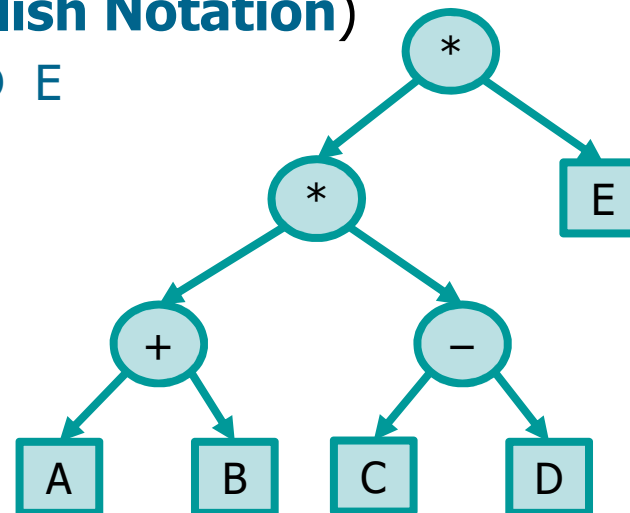
Example

$[(A + B) * (C - D)] * E$

- ❖ A pre-order visit returns the expression in the seldom used **prefix** form (**Polish Notation**)

➤ * * + A B - C D E

Brackets are no more needed



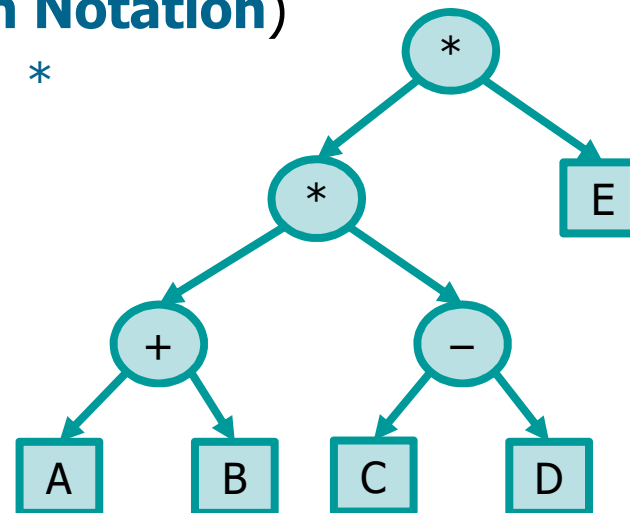
Example

$[(A + B) * (C - D)] * E$

- ❖ A post-order visit returns the expression in **postfix** form (**Reverse Polish Notation**)

➤ $A B + C D - * E *$

Brackets are no more needed



Exam: 29 January 2018

Exercise

- ❖ Convert the following expressions from in-fix to post-fix and pre-fix notations

ECE students (10 credits)

$$(A - B) / \{ (C / D) + [(D / (E - F))] \}$$

$$(A - B) / \{ (C / D) + [(D / (E - F)) * G] \}$$

CE students (12 credits)

A parser for the prefix form

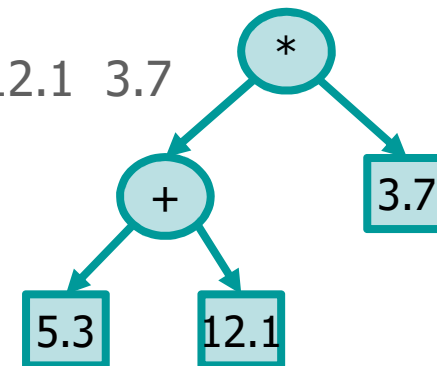
- ❖ The following grammar specifies the prefix form (Polish notation)

```
<exp> = <operand> | <op> <exp> <exp>  
<operand> = float  
<op> = + | * | - | /
```

- Example

- $(5.3 + 12.1) * 3.7 \rightarrow * + 5.3 12.1 3.7$

- ❖ Write a recursive program to implement this grammar



Implementation

```
int main(int argc, char *argv[]) {
    float result;
    int pos=0;

    if (argc < 2) {
        fprintf(stderr, "Error: missing parameter.\n");
        fprintf(stderr, "Run as: %s prefix_expression\n",
            argv[0]);
        return 1;
    }

    result = eval_r(argv[1], &pos);
    fprintf(stdout, "Result = %.2f\n", result);
    return EXIT_SUCCESS;
}
```

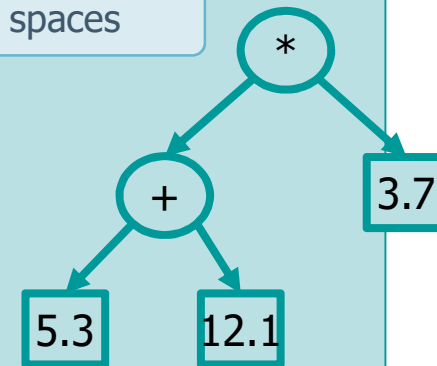
Implementation

Expression

```
float eval_r (char *expr, int *pos_ptr) {  
    float left, right, result;  
    char operator;  
    int k = *pos_ptr;  
    while (isspace(expr[k])) {  
        k++;  
    }  
    if (expr[k]=='+' || expr[k]=='*' ||  
        expr[k]=='-' || expr[k]=='/') {  
        operator = expr[k++];  
        left = eval_r(expr, &k);  
        right = eval_r(expr, &k);  
        switch (operator) {  
            case '+': result = left+right; break;  
            case '*': result = left*right; break;  
            case '-': result = left-right; break;  
            case '/': result = left/right; break;  
        }  
    }  
}
```

Parsing index

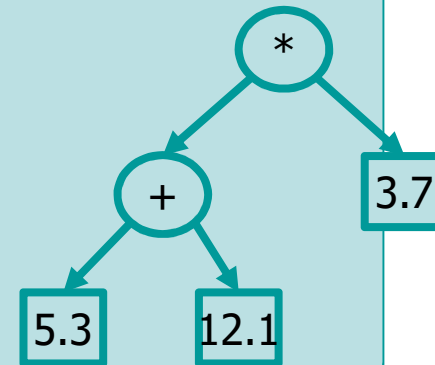
Skip spaces



Implementation

```
} else {  
    sscanf(&expr[k], "%f", &result);  
    while (isdigit(expr[k]) || expr[k]=='.') {  
        k++;  
    }  
}  
  
*pos_ptr = k;  
return result;  
}
```

Terminal case:
A real value



A parser for the postfix form

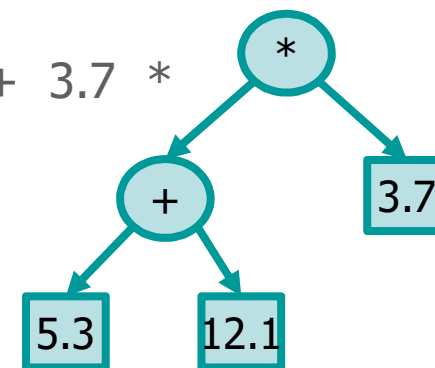
- ❖ The following grammar specifies the postfix form (Polish notation)

```
<exp> = <operand> | <exp> <exp> <op>  
<operand> = float  
<op> = + | * | - | /
```

- Example

- $(5.3 + 12.1) * 3.7 \rightarrow 5.3 \ 12.1 \ + \ 3.7 \ *$

- ❖ Write a program to implement this grammar



Implementation

```
#include <stdio.h>
#include <string.h>
#include <ctype.h>

#include "util.h"
#include "stackPublic.h"

int main(int argc, char *argv[]) {
    float result;
    int left, right, length, k=0;
    stack_t *sp=NULL;
    char *expr;

    util_check_m(argc>=2, "missing parameter.");
    expr = argv[1];
    length = strlen(expr);
    sp = stack_init(length);
```

ADT for utility functions

ADT for the stack

Implementation

```
while (k < length) {
    if (isdigit(expr[k])) {
        sscanf(&expr[k], "%f", &result);
        stack_push(sp, (void *)result);
        while (isdigit(expr[k]) || expr[k]=='.') {
            k++;
        }
    } else if (expr[k]=='+' || expr[k]=='*' ||
               expr[k]=='-' || expr[k]=='/') {
        stack_pop(sp, (void **)&right);
        stack_pop(sp, (void **)&left);
        switch (expr[k]) {
            case '+': result = left+right; break;
            case '*': result = left*right; break;
            case '-': result = left-right; break;
            case '/': result = left/right; break;
        }
        stack_push(sp, (void *)result);
    }
    k++;
}
```

Skip float

Implementation

```
stack_pop(sp, (void **)&result);  
fprintf(stdout, "Result = %ld\n", result);  
stack_dispose(sp, NULL);  
  
return EXIT_SUCCESS;  
}
```